Dual Encryption by Random Segmentation and Random Re-Arrangement (RSRA) using Two Dimensional Array

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Abstract—**Cryptography is a method of storing and transmitting data in a particular form so that only those for whom it is intended can read and process it. Cryptography includes techniques such as microdots, merging words with images, and other ways to hide information in storage or transit. However, in today's computer-centric world, cryptography is most often associated with scrambling plaintext (ordinary text, sometimes referred to as cleartext) into ciphertext (a process called encryption), then back again (known as decryption). Individuals who practice this field are known as cryptographers.**

To make a secure and successful transaction of message, the security of cipher text is must before transmission over communication channel. This paper propose a new encryption technique with the help of random segmentation and random re-arrangement of cipher text using two dimensional array before transmission over communication channel so that the transmitted cipher text is more complex to decrypt by any attacker.

Index Terms—**Crptography, Symmetric encryption, Asymmetric encryption, Cryptanalysis, Cipher text, AES, DES, Dual encryption, RSA, ECC, *Adaptive chosen-plaintext, Related-key attack, Private Key, Public Key, Secret Key.***

I. INTRODUCTION

Until modern times, cryptography referred almost exclusively to *encryption*, which is the process of converting ordinary information (called plaintext) into unintelligible text (called ciphertext). Decryption is the reverse, in other words, moving from the unintelligible ciphertext back to plaintext. A *cipher* (or *cypher*) is a pair of algorithms that create the encryption and the reversing decryption. The detailed operation of a cipher is controlled both by the algorithm and in each instance by a "key". The key is a secret (ideally known only to the communicants), usually a short string of characters, which is needed to decrypt the ciphertext. Formally, a "cryptosystem" is the ordered list of elements of finite possible plaintexts, finite possible cyphertexts, finite possible keys, and the encryption and decryption algorithms which correspond to each key. Keys are important both formally and in actual practice, as ciphers without variable keys can be trivially broken with only the knowledge of the cipher used and are therefore useless (or even counter-productive) for most purposes. Historically, ciphers were often used directly for encryption or decryption without additional procedures such as authentication or integrity checks.

II. LITERATURE REVIEW

Modern encryption methods can be divided by two criteria: by type of key used, and by type of input data.
1. By type of key used ciphers are divided into:
   1.1 Symmetric key algorithms (Private-key cryptography)
In symmetric systems the same key (the secret key) is used to encrypt and decrypt a message. Data manipulation in symmetric systems is faster than asymmetric systems as they generally use shorter key lengths. Symmetric models include the commonly used AES (Advanced Encryption Standard) which replaced the older DES (Data Encryption Standard).
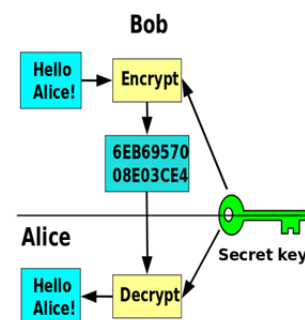
(Figure 1.1 : Symmetric-key cryptography, where a single key is used for encryption and decryption)

1.2 Asymmetric key algorithms (Public-key cryptography)

Asymmetric systems use a public key to encrypt a message and a private key to decrypt it. Use of asymmetric systems enhances the security of communication. Examples of asymmetric systems include RSA (Rivest-Shamir-Adleman), and ECC (Elliptic Curve Cryptography).
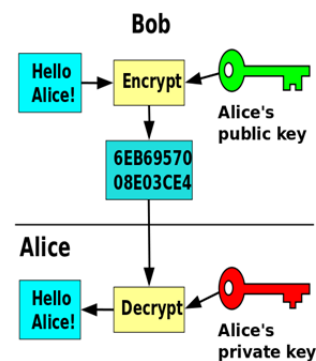
(Figure 1.2 : Public-key cryptography, where different keys are used for encryption and decryption)

2. Ciphers can be distinguished into two types by the type of input data:

2.1 Block ciphers
which encrypt block of data of fixed size, and
2.2 Stream ciphers
which encrypt continuous streams of data
Although there are various algorithm to encrypt any plain text to cipher text but there are also various legal and illegal activities around the world in which any one try to steal our encrypted message over the communication channel and try to decode it to get the access of plain text.

Cryptanalysis is the study of analyzing information systems in order to study the hidden aspects of the systems. Cryptanalysis is used to breach cryptographic security systems and gain access to the contents of encrypted messages, even if the cryptographic key is unknown. The goal of the *cryptanalyst* is to gain as much information as possible about the original, unencrypted data.

Attacks can be classified based on what type of information the attacker has available:
- *Ciphertext-only*: the cryptanalyst has access only to a collection of ciphertexts or codetexts.
- *Known-plaintext*: the attacker has a set of ciphertexts to which he knows the corresponding plaintext.
- *Chosen-plaintext* (*chosen-ciphertext*): the attacker can obtain the ciphertexts (plaintexts) corresponding to an arbitrary set of plaintexts (ciphertexts) of his own choosing.
- *Adaptive chosen-plaintext*: like a chosen-plaintext attack, except the attacker can choose subsequent plaintexts based on information learned from previous encryptions. Similarly *Adaptive chosen ciphertext attack*.
- *Related-key attack*: Like a chosen-plaintext attack, except the attacker can obtain ciphertexts encrypted under two different keys. The keys are unknown, but the relationship between them is known; for example, two keys that differ in the one bit.

III. PROPOSED SYSTEM

So now we can see that there can be any type of attack on our cipher text over the communication channel by which any attacker can get the access to our original message by trying any of attacking method. So it is very necessary that we apply a method, which will make our cipher text more secure and unbreakable over the communication channel. In this proposed method 1st we convert plain text to middle level cipher text by using any of encryption technique (Symmetric or Asymmetric) and 2nd we randomly rearrange this middle level cipher text by splitting the cipher text in random number of parts and then positioning each part in random sequence and by storing all these information in a 2 dimensional array.

Following are the steps in the proposed method :

Step 1 : Convert the Plain text to Cipher Text by using any of encryption technique (*Symmetric key algorithm in this example*)
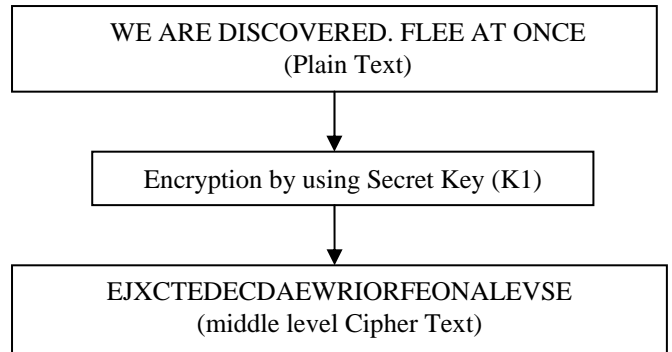
Figure 2.1

Step 2 : Split the cipher text in random number of blocks with random length by using following algorithm

```
n = number of alphabets in cipher text.
c= cipher text.
sc=0. (starting index)
ec=0. (end index)
i=0;
Array a[][];

while (ec<=n)

start

if sc=0
start
ec=random(0,n-2);
a[0][0]=sc;
a[0][1]=ec;
sc=ec+1;
i++;
end

else
start
a[i][0]=sc;
ec=random(sc,n-1);
a[i][1]=ec;
sc=ec+1;
i++;
end

end while loop
```

For example, if we apply this process on above cipher text

```
0  1  2  3  4  5  6  7  8  9 10 11 12 13 14 15 16 17 18
E  J  X  C  T  E  D  E  C  D  A  E  W  R  I  O  R  F  E
   _____       _____    ____    ____    _____
      1              2          3       4         5
19 20 21 22 23 24 25 26
O  N  A  L  E  V  S  E
   _____  ____  _  ____
      6       7    8    9
```

n = 27
c= EJXCTEDECDAEWRIORFEONALEVSE
sc=0. (starting index)
ec=0. (end index)
i=0;

in 1st while loop:
sc=0;
ec=5 (let's consider)=random(1,25);
a[0][0]=0;
a[0][1]=5;

in 2nd while loop:
sc=5+1=6;
ec=8 (let's consider)=random(6,26);
a[1][0]=6;
a[1][1]=8;

and so on...

By following above algorithm, after end of while loop the array 'a' will look as follow:

sc	ec	Seq.
0	5	
6	8	
9	12	
13	14	
15	18	
19	21	
22	23	
24	24	
25	26	

Table 1.1

Step 3 : Now assign random sequence no to each row of array a

For Example :

sc	ec	Seq.
0	5	6
6	8	2
9	12	5
13	14	9
15	18	1
19	21	4
22	23	8
24	24	3
25	26	7

Table 1.2

Step 4 : Now arrange each segment of cipher text according to its randomly assigned sequence number.

O R F E D E C V O N A D A E W E J X C T
E S E L E R I

(Final Cipher Text)

Now this final ciphertext can be transmitted over communication channel. The Array 'a' is converted into a linear arrangement of each row of array 'a' according to randomly assigned sequence number and will be provided to receiver by any secret method as a key to decrypt the final cipher text to middle level text. Again this middle level cipher text will be converted to plain text by using Secret Key K1.

In above example the information of Array 'a' will be arranged as follow :

K2   15 18 1 6 8 2 24 24 3 19 21 4 9 12 5 0 5 6 25 26 7 22 23 8 13 14 9

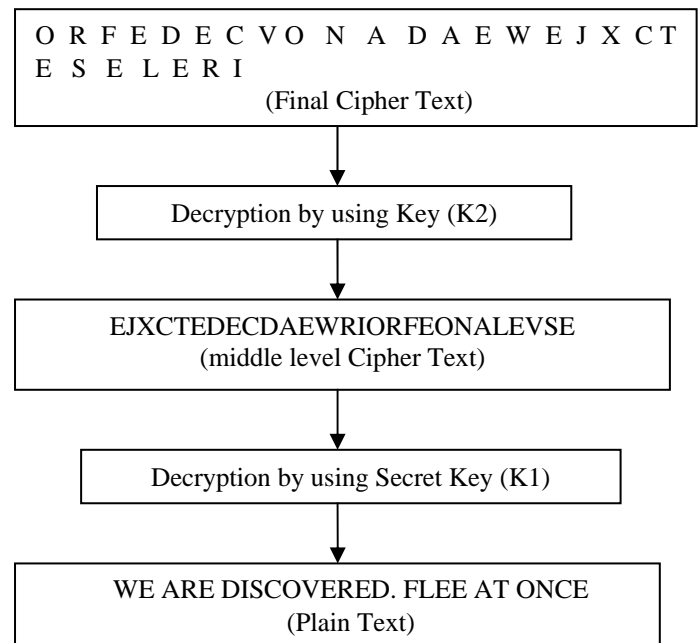The reverse process will be applied for decryption of above cipher text to the plain text.

O R F E D E C V O N A D A E W E J X C T
E S E L E R I
(Final Cipher Text)

↓

Decryption by using Key (K2)

↓

EJXCTEDECDAEWRIORFEONALEVSE
(middle level Cipher Text)

↓

Decryption by using Secret Key (K1)

↓

WE ARE DISCOVERED. FLEE AT ONCE
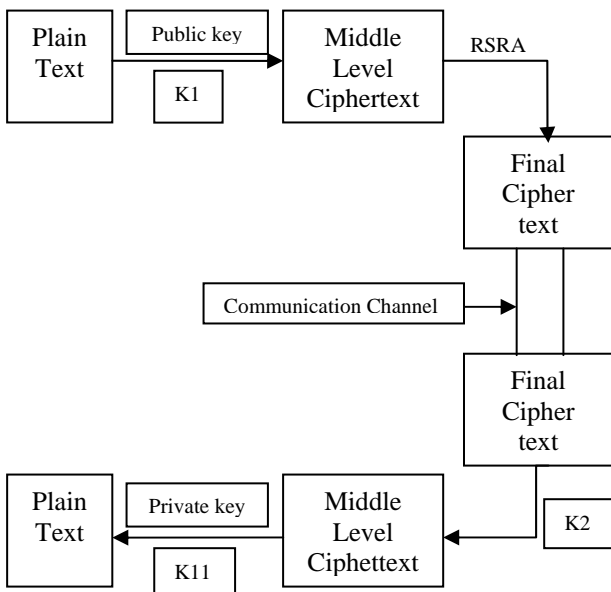(Plain Text)

Figure 2.2

This proposed method can be applied on both Symmetric and Asymmetric encryption algorithm as follow :

A. **For Symmetric encryption algorithms:**
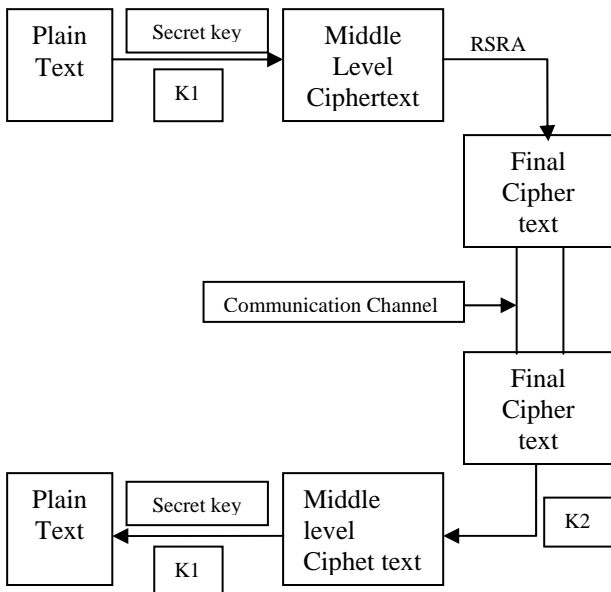
Figure 2.3

B. **For Asymmetric encryption algorithms:**

Figure 2.4

IV. CONCLUSION

Hence as described in above proposed method, the cipher text can be made more complex with the help of random segmentation and random re-arrangement of cipher text using two dimensional array before transmission over communication channel so that the transmitted cipher text can not be decrypted by any attacker.

V. ACKNOWLEDGMENT

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Er. Amit Kumar received his Bachelor of Technology (B.Tech) degree in Computer Science and engineering in 2010 from Subharti Institute of Technology and Engineering (SITE) Meerut, India. He has around six year industrial working experience as a project leader. He is currently working towards Master of Technology (M.Tech) degree in Cyber Security at Subharti Institute of Technology and Engineering(SITE), Swami Vivekanad Subharti University Meerut, U.P. (India) . His area of research includes Software Engineering(Secured Web ERP Development), Cryptography, Cryptanalysis and Algorithms.